

# A NOTE ON THE YONEDA LEMMA

ALEX MASSARENTI

ABSTRACT. We state, without proofs, the Yoneda lemma, and the 2-Yoneda lemma. Then we prove that a functor is representable if and only if it admits a universal object using the Yoneda lemma, and that a fibered category is representable if and only if it is fibered in groupoids using the 2-Yoneda lemma.

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## 1. THE YONEDA LEMMA

Let  $\mathcal{C}$  be a category. The functors from  $\mathcal{C}^{op}$  to  $(Set)$  can be thought as a category  $Hom(\mathcal{C}^{op}, (Set))$ , in which the arrows are the natural transformations. To any object  $X \in \mathcal{C}$  we can associate a functor  $h_X : \mathcal{C}^{op} \rightarrow (Set)$ , which maps an object  $Y \in \mathcal{C}$  in  $Hom_{\mathcal{C}}(Y, X)$ , and a morphism  $f : Y \rightarrow Z$  in the map

$$h_X f : Hom_{\mathcal{C}}(Z, X) \rightarrow Hom_{\mathcal{C}}(Y, X), \alpha \mapsto \alpha \circ f.$$

Now let now  $\rho : X \rightarrow Y$  be an arrow in  $\mathcal{C}$ . For any object  $Z \in \mathcal{C}$  we get a function  $h_f Z : h_X Z \rightarrow h_Y Z$ ,  $\alpha \mapsto \rho \circ \alpha$ . This defines a morphism  $h_f : h_X \rightarrow h_Y$ , i.e. for any arrow  $\beta : Z \rightarrow W$  in  $\mathcal{C}$  the following diagram commute

$$\begin{array}{ccc} h_X Z & \xrightarrow{h_f Z} & h_Y Z \\ h_X \beta \downarrow & & \downarrow h_Y \beta \\ h_X W & \xrightarrow{h_f W} & h_Y W \end{array}$$

In this way we have defined a functor

$$\mathcal{H} : \mathcal{C}^{op} \rightarrow Hom(\mathcal{C}^{op}, (Set)), X \rightarrow h_X, f \rightarrow h_f.$$

**Lemma 1.1.** (*Weak Yoneda Lemma*) *The functor  $\mathcal{H}$  is fully faithful, that is the function*

$$Hom(X, Y) \rightarrow Hom(h_X, h_Y), f \mapsto h_f$$

*is bijective.*

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The functor  $\mathcal{H}$  is not an equivalence of categories because it is not essentially surjective. This means that there are functors  $\mathcal{C}^{op} \rightarrow (Set)$  not isomorphic to a functor of the form  $h_X$ . The lemma says that the category  $\mathcal{C}$  can be embedded in the category  $Hom(\mathcal{C}^{op}, (Set))$ .

**Definition 1.2.** A functor  $F : \mathcal{C}^{op} \rightarrow (Set)$  is said to be representable if it is isomorphic to a functor of the form  $h_X$  for some  $X \in \mathcal{C}^{op}$ .

If we restrict to the full subcategory of  $Hom(\mathcal{C}^{op}, (Set))$  of representable functor then  $\mathcal{H}$  is an equivalence. Note that if  $F \cong h_X$  and  $F \cong h_Y$  then the isomorphism between  $h_X$  and  $h_Y$  comes from a unique isomorphism from  $X$  and  $Y$ , by the weak Yoneda lemma. This means that two objects representing the same functor are canonically isomorphic.

Our aim is to reformulate the definition of representable functor using a more general version of Yoneda's lemma. Let  $X \in \mathcal{C}$  be an object and let  $F : \mathcal{C}^{op} \rightarrow (Set)$  be a functor. Given a natural transformation  $\alpha : h_X \rightarrow F$  we can consider the map  $\alpha_X : h_X(X) = Hom(X, X) \rightarrow F(X)$ , and the element  $\xi = \alpha_X(Id_X)$ . We get a function

$$Hom(h_X, F) \rightarrow F(X), \alpha \mapsto \alpha_X(Id_X) = \xi.$$

Conversely fixed an element  $\xi \in F(X)$ , we define a natural transformation  $\alpha : h_X \rightarrow F$  as follows. Let  $Y \in \mathcal{C}$  be an object, and let  $f \in h_X(Y) = Hom(Y, X)$  be an arrow, the arrow  $f$  induces a function  $F(f) : F(X) \rightarrow F(Y)$ . We define  $\alpha_Y : h_X(Y) \rightarrow F(Y)$  mapping  $f \mapsto F(f)(\xi)$ . We have defined a function

$$F(X) \rightarrow Hom(h_X, F).$$

**Lemma 1.3.** (*Yoneda Lemma*) *The two function above define a bijective correspondence*

$$Hom(h_X, F) \xrightarrow{\cong} F(X)$$

Note that for  $F = h_Y$  we get a bijective correspondence  $Hom(h_X, h_Y) \xrightarrow{\cong} h_Y(X) = Hom(X, Y)$ , that is the weak Yoneda lemma.

**Definition 1.4.** Let  $F : \mathcal{C}^{op} \rightarrow (Set)$  be a functor. An *universal object* for  $F$  is a pair  $(X, \xi)$  where  $X \in \mathcal{C}$  and  $\xi \in F(X)$  such that for each  $Y \in \mathcal{C}$  and each  $\beta \in F(Y)$ , there is a unique arrow  $f : Y \rightarrow X$  such that  $\beta = F(f)(\xi)$ .

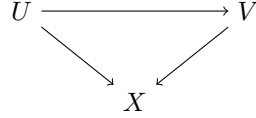
Recall that for any  $\xi \in F(X)$  we defined a morphism  $\alpha : h_X \rightarrow F$ , in other words  $(X, \xi)$  is an universal object if and only if the morphism  $\alpha$  defined by  $\xi \in F(X)$  is an isomorphism. By Yoneda lemma any natural transformation  $h_X \rightarrow F$  is induced by a  $\xi \in F(X)$ . So we get the following proposition.

**Proposition 1.5.** *A functor  $F : \mathcal{C}^{op} \rightarrow (Set)$  is representable if and only if it has an universal object.*

Clearly if  $(X, \xi)$  is an universal object for  $F$  then  $X$  represents  $F$ . The Yoneda lemma establish an equivalence of categories between  $\mathcal{C}$  and the full subcategory of  $Hom(\mathcal{C}^{op}, (Set))$  defined by representable functors. In particular if  $\mathcal{C} = (Sch/k)$  is the category of schemes over the field  $k$ , we get that a scheme  $X$  is determined by its functor of points  $h_X : (Sch/k)^{op} \rightarrow (Set)$ .

2. THE 2-YONEDA LEMMA

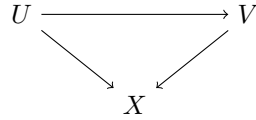
Let  $\mathcal{C}$  be a category and let  $X \in \mathcal{C}$  be an object. We can construct the *Comma category*  $(\mathcal{C}/X)$  as follows. The objects of  $(\mathcal{C}/X)$  are morphism  $U \rightarrow X$  in  $\mathcal{C}$ , and the arrows between two objects  $U \rightarrow X$  and  $V \rightarrow X$  are the arrows  $U \rightarrow V$  in  $\mathcal{C}$  such that the following diagram commutes



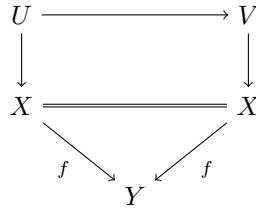
The fibered category  $(\mathcal{C}/X) \rightarrow \mathcal{C}$ ,  $(U \rightarrow X) \rightarrow U$  is the category fibered in sets associated to the functor  $h_X : \mathcal{C}^{op} \rightarrow (Set)$ . Recall that we have an embedding  $\mathcal{C} \rightarrow Hom(\mathcal{C}^{op}, (Set))$ ,  $X \mapsto h_X$ , and by composing we get an embedding

$$\mathcal{C} \rightarrow Hom(\mathcal{C}^{op}, (Set)) \rightarrow (FC/\mathcal{C}), X \mapsto h_X \mapsto ((\mathcal{C}/X) \rightarrow \mathcal{C}),$$

of  $\mathcal{C}$  in  $(FC/\mathcal{C})$ , the 2-category of fibered categories over  $\mathcal{C}$ . Note that a morphism  $f : X \rightarrow Y$  in  $\mathcal{C}$  goes to the morphism  $(\mathcal{C}/f) : (\mathcal{C}/X) \rightarrow (\mathcal{C}/Y)$  of fibered categories defined as follows. Let  $U \rightarrow X$  be an object in  $(\mathcal{C}/X)$  then  $(\mathcal{C}/f)(U \rightarrow X) = (U \rightarrow X \rightarrow Y)$  i.e. a morphism goes in its composition with  $f$ . The functor  $(\mathcal{C}/f)$  sends an arrow



in  $(\mathcal{C}/X)$ , in the commutative diagram



**Lemma 2.1.** (*Weak 2-Yoneda Lemma*) *The function that sends  $f : X \rightarrow Y$  in the morphism  $(\mathcal{C}/f) : (\mathcal{C}/X) \rightarrow (\mathcal{C}/Y)$  states a bijection between  $Hom_{\mathcal{C}}(X, Y)$  and  $Hom((\mathcal{C}/X), (\mathcal{C}/Y))$ .*

**Definition 2.2.** A fibered category over  $\mathcal{C}$  is representable if it is equivalent to a category of the form  $(\mathcal{C}/X)$ .

Let  $F$  be a category fibered over  $\mathcal{C}^{op}$ , and let  $X$  be an object of  $\mathcal{C}^{op}$ . Let  $f \in Hom_{\mathcal{C}}((\mathcal{C}/X), F)$  be a morphism of fibered categories. Then we have  $f_X : (\mathcal{C}/X)(X) \rightarrow F(X)$  and we can consider  $f(Id_X) \in F(X)$ . Furthermore if  $\alpha : f \rightarrow g$  is a base preserving natural transformation between  $f, g : (\mathcal{C}/X)raF$ , we get an arrow  $\alpha_{Id_X} : f(Id_X) \rightarrow g(Id_X)$ . This defines a functor

$$Hom_{\mathcal{C}}((\mathcal{C}/X), F) \rightarrow F(X).$$

Conversely, if  $\xi \in F(X)$  is an object we can construct a functor  $f_\xi : (\mathcal{C}/X) \rightarrow F$  as follows. For any  $\psi : U \rightarrow X$  in  $(\mathcal{C}/X)$  we define  $f_\xi(\psi) = \psi^*\xi \in F(U)$ . If

$$\begin{array}{ccc} U & \xrightarrow{g} & V \\ & \searrow \alpha & \swarrow \beta \\ & & X \end{array}$$

is an arrow in  $(\mathcal{C}/X)$  we associate the unique arrow  $\rho : \alpha^*\xi \rightarrow \beta^*\xi$  in  $F$  such that the diagram

$$\begin{array}{ccccc} \alpha^*\xi & & & & \\ & \searrow \rho & & & \\ & & \beta^*\xi & \longrightarrow & \xi \\ & & \downarrow \alpha & & \downarrow \\ U & & V & \xrightarrow{\beta} & X \\ & \searrow g & & & \\ & & & & \end{array}$$

commutes.

**Lemma 2.3.** (*2-Yoneda Lemma*) *The functors above define an equivalence of categories*

$$\text{Hom}_{\mathcal{C}}((\mathcal{C}/X), F) \cong F(X).$$

Then a morphism  $(\mathcal{C}/X) \rightarrow F$  corresponds to an object  $\xi \in F(X)$ , which in turn defines the functor  $f_\xi : (\mathcal{C}/X) \rightarrow F$  defined above, this is isomorphic to the original functor  $f$ . Now the fibered category  $F$  is representable if and only if  $f_\xi : (\mathcal{C}/X) \rightarrow F$  is an equivalence for some  $X \in \mathcal{C}$  and  $\xi \in F(X)$ . Recall that  $f_\xi$  is an equivalence if and only if for each object  $Y \in \mathcal{C}^{op}$  the functor

$$f_\xi(Y) : (\mathcal{C}/X)(Y) = \text{Hom}_{\mathcal{C}}(Y, X) \rightarrow F(Y), g \mapsto g^*\xi,$$

is an equivalence of categories. Since  $\text{Hom}_{\mathcal{C}}(Y, X)$  is a set this means that  $F(Y)$  is a groupoid. Thanks to the 2-Yoneda lemma we get another characterization for representable fibered categories.

**Proposition 2.4.** *A fibered category  $F$  over  $\mathcal{C}$  is representable if and only if it is fibered in groupoids, and there exist an object  $Y \in \mathcal{C}$ , and an object  $\xi \in F(Y)$ , such that for any object  $\nu$  of  $F$  there exists a unique arrow  $\nu \rightarrow \xi$  in  $F$ .*

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SISSA, VIA BONOMEA 265, 34136 TRIESTE, ITALY  
*E-mail address:* alex.massarenti@sissa.it