

# Inductive Logic Programming

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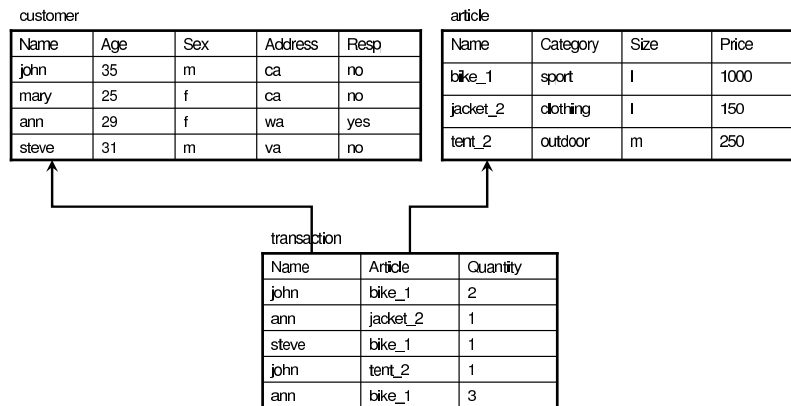
# Targeted Mailing

customer

Name	Age	Sex	Address	Resp
john	35	m	ca	no
mary	25	f	ca	yes
ann	29	f	wa	no
steve	31	m	va	no

If Age<30 and Address=ca then Resp=yes

# Multi Table



The customer will respond if she/he has bought an item of category clothing

# Join

customer >< transaction >< article

Name	Age	Sex	Address	Article	Quantity	Category	Size	Price	Resp
john	35	m	ca	bike_1	2	sport	l	1000	no
john	35	m	ca	tent_2	1	outdoor	m	250	no
mary	25	f	ca						no
ann	29	f	wa	jacket_2	1	clothing	l	150	yes
ann	29	f	wa	bike_1	3	sport	l	1000	yes
steve	31	m	va	bike_1	2	sport	l	1000	no

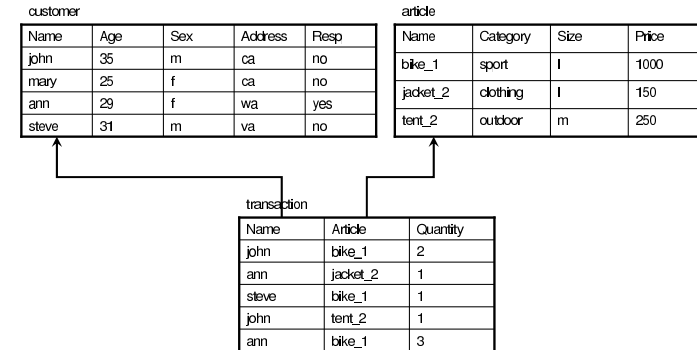
# Replicate

Replicate attributes

Name	Age	Sex	Address	Article1	Quantity1	Category1	Size1	Price1
john	35	m	ca	bike_1	2	sport	l	1000
mary	25	f	ca					
ann	29	f	wa	jacket_2	1	clothing	l	150
steve	31	m	va	bike_1	2	sport	l	1000

Article2	Quantity2	Category2	Size2	Price2	Resp
tent_2	1	outdoor	m	250	no
					no
bike_1	3	sport	l	1000	yes
					no

# Logic



$respond(Customer) \leftarrow$   
 $transaction(Customer, Article, \_Quantity),$   
 $article(Article, Category, \_Size, \_Price),$   
 $Category = clothing.$

# Outline of the Talk

- Predictive ILP
  - Learning from entailment
    - Bottom-up systems: Golem
    - Top-down systems: FOIL, Progol
  - Learning from interpretations
    - ICL, Tilde
- Descriptive ILP
  - Claudien
- Probabilistic ILP
  - ALLPAD
- Applications

# Predictive ILP

- Aim:
  - classifying instances of the domain, i.e.
  - predicting the class
- Two settings:
  - Learning from entailment
  - Learning from interpretations

## Learning from Entailment

- Given
  - A set of positive example  $E^+$
  - A set of negative examples  $E^-$
  - A background knowledge  $B$
  - A space of possible programs  $\mathcal{H}$
- Find a program  $P \in \mathcal{H}$  such that
  - $\forall e^+ \in E^+, P \cup B \models e^+$  (completeness)
  - $\forall e^- \in E^-, P \cup B \not\models e^-$  (consistency)

## Mailing Example

- Positive examples  $E^+ = \{respond(ann)\}$
- Negative examples  
 $E^- = \{respond(john), respond(mary), respond(steve)\}$
- Background  $B =$  facts for relations *customer*,  
*transaction* and *article*  
 $customer(john, 35, m, ca)$ .  
 $customer(mary, 25, f, ca)$ .  
 $customer(ann, 29, f, wa)$ . . . .  
 $transaction(john, bike\_1, 2)$ .  
 $transaction(ann, jacket\_2, 1)$ . . . .  
 $article(bike\_1, sport, l, 1000)$ .  
 $article(jacket\_2, clothing, 1, 150)$ . . . .

## Mailing Example

- Space of programs  $\mathcal{H}$ : programs containing clauses with
  - in the head  $respond(Customer)$
  - in the body a conjunction of literals from the set  
 $\{customer(Customer, Age, Sex, Address),$   
 $transaction(Customer, Article, Quantity),$   
 $article(Article, Category, Price),$   
 $Age = constant, Sex = constant, \dots\}$

## Definitions

- $covers(P, e) = true$  if  $B \cup P \models e$
- $covers(P, E) = \{e \in E \mid covers(P, e) = true\}$
- A theory  $P$  is more general than  $Q$  if  
 $covers(P, U) \supseteq covers(Q, U)$
- If  $B \cup P \models Q$  then  $P$  is more general than  $Q$
- A clause  $C$  is more general than  $D$  if  
 $covers(\{C\}, U) \supseteq covers(\{D\}, U)$
- If  $B, C \models D$  then  $C$  is more general than  $D$
- If a clause covers an example, all of its generalizations will (*covers* is antimonotonic)
- If a clause does not cover an example, none of its specializations will

## Theta Subsumption

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- $C$   $\theta$ -subsumes  $D$  ( $C \geq D$ ) if there exists a substitution  $\theta$  such that  $C\theta \subseteq D$  [Plotkin 70]
- $C \geq D \Rightarrow C \models D \Rightarrow B, C \models D \Rightarrow C$  is more general than  $D$
- $C \models D \not\Rightarrow C \geq D$

## Examples of Theta Subsumption

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- $C1 = grandfather(X, Y) \leftarrow father(X, Z)$
- $C2 = grandfather(X, Y) \leftarrow father(X, Z), parent(Z, Y)$
- $C3 = grandfather(john, steve) \leftarrow father(john, mary), parent(mary, steve)$
- $C1 \geq C2$  with  $\theta = \emptyset$
- $C1 \geq C3$  with  $\theta = \{X/john, Y/steve, Z/mary\}$
- $C2 \geq C3$  with  $\theta = \{X/john, Y/steve, Z/mary\}$

## In Practice

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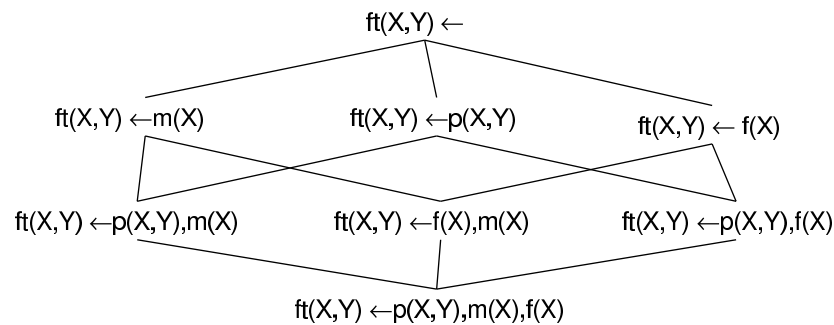
- Logical consequence is undecidable
- Coverage test: SLD or SLDNF resolution
- Generality order:
  - $\theta$ -subsumption because it is decidable (even if NP-complete)

## Properties of Theta Subsumption

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- $\theta$ -subsumption induces a lattice in the space of clauses
- Every set of clauses has a least upper bound (lub) and a greatest lower bound (glb)
- This is not true for the generality relation based on logical consequence

## Lattice



## Least General Generalization

- $lgg(C, D)$  = least upper bound in the  $\theta$ -subsumption order
- An algorithm exists which has complexity  $O(s^2)$  where  $s$  is the size of the clauses
- Example:

$C = \text{father}(\text{john}, \text{mary}) \leftarrow \text{parent}(\text{john}, \text{mary}), \text{male}(\text{john})$   
 $D = \text{father}(\text{david}, \text{steve}) \leftarrow \text{parent}(\text{david}, \text{steve}), \text{male}(\text{david})$   
 $lgg(C, D) = \text{father}(X, Y) \leftarrow \text{parent}(X, Y), \text{male}(X)$

- For a set of  $n$  clauses the complexity is  $O(s^n)$

## Relative Subsumption

- $\theta$  subsumption does not take into account background knowledge
- $C \geq D \Leftrightarrow \models \forall (C\theta \rightarrow D)$
- Relative Subsumption [Plotkin 71]:  $C$   $\theta$  subsume  $D$  relative to background  $B$  ( $C \geq_B D$ ) if there exists a substitution  $\theta$  such that  $B \models \forall (C\theta \rightarrow D)$

## Relative Least General Generalization

- Relative Least General Generalization (rlgg): lgg with respect to relative subsumption.
- It does not exist in the general case of  $B$  a set of Horn clauses
- It exists in the case that  $B$  is a set of ground atoms and can be computed in this way:
- $rlgg((H1 \leftarrow B1), (H2 \leftarrow B2)) = lgg((H1 \leftarrow B1, B), (H2 \leftarrow B2, B))$
- If  $B$  is a set of clauses, we can compute the  $h$ -easy model

## Relative Least General Generalization

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- Example

$C1 = \text{father}(\text{john}, \text{mary})$

$C2 = \text{father}(\text{david}, \text{steve})$

$B = \{\text{parent}(\text{john}, \text{mary}), \text{parent}(\text{david}, \text{steve}),$   
 $\text{parent}(\text{kathy}, \text{ellen}), \text{female}(\text{kathy}),$   
 $\text{male}(\text{john}), \text{male}(\text{david})\}$

$\text{rlgg}(C1, C2) = \text{father}(X, Y) \leftarrow \text{parent}(X, Y), \text{male}(X)$

## Outline of the Talk

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    - **Bottom-up systems: Golem**
    - Top-down systems: FOIL, Progol
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    - ICL, Tilde
- Descriptive ILP
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## Bottom-up Systems

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- Covering loop
- Search for a clause from specific to general

**Learn**( $E, B$ )

$P := \emptyset$

repeat /\* covering loop \*/

$C := \text{GenerateClauseBottomUp}(E, B)$

$P := P \cup \{C\}$

    Remove from  $E$  the positive examples covered by  $P$

until Sufficiency criterion

return  $P$

## Golem [Muggleton, Feng 90]

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- Bottom-up system
- Generalization by means of rlgg
- Sufficiency criterion:  $E^+ = \emptyset$

# Golem

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## **GolemGenerateClause**( $E, B$ )

select randomly some couples of examples from  $E^+$

compute their rlgg

let  $C$  be the rlgg that covers most positive examples

while covering no negative

repeat

randomly select some examples from  $E^+$

compute the rlgg between  $C$  and each selected example

let  $C$  be the rlgg that covers most positive examples

while covering no negative

remove from  $E^+$  the examples covered by  $C$

while the set of examples covered by  $C$  increases

remove literals from the body of  $C$  until  $C$  covers

some negative examples

return  $C$

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- 

# Top-down Systems

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- Covering loop as bottom-up systems
  - Search for a clause from general to specific
- 

# Top-down Systems

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## **GenerateClauseTopDown**( $E, B$ )

$Beam := \{p(X) \leftarrow true\}$

$BestClause := null$

repeat /\* specialization loop \*/

Remove the first clause  $C$  of  $Beam$

compute  $\rho(C)$

score all the refinements

update  $BestClause$

add all the refinements to the beam

order the beam according to the score

remove the last clauses that exceed the dimension  $d$

until the Necessity criterion is satisfied

return  $BestClause$

---

## Typical Stopping Criteria

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- Sufficiency criteria:
  - $E^+ = \emptyset$
  - `GenerateClauseTopDown` returns *null*
  - a disjunction of the above
- Necessity criteria
  - the number of negative examples covered by *BestClause* is 0
  - the number of covered negative examples covered by *BestClause* is below a threshold
  - *Beam* is empty
  - a disjunction of the above

## Refinement Operator

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- $\rho(C) = \{D \mid D \in L, C \geq D\}$
- where  $L$  is the space of possible clauses
- A refinement operator usually generates only minimal specializations
- A typical refinement operator applies two syntactic operations to a clause
  - applies a substitution to the clause
  - adds a literal to the body

## Heuristic Functions

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- Notation:  $n^+(C), n^-(C)$  number of positive and negative examples covered by clause  $C$
- $n(C) = n^+(C) + n^-(C)$
- Accuracy:  $Acc = p(+|C)$  (more accurately Precision),  $p(+|C)$  can be estimated by
  - relative frequency:  $p(+|C) = \frac{n^+(C)}{n(C)}$
  - m-estimate:  $p(+|C) = \frac{n^+(C) + mp(+)}{n(C) + m}$ , where  $p(+)=n^+/n$
  - Laplace: m-estimate with  $m = 2, p = 0.5$   
 $p(+|C) = \frac{n^+(C) + 1}{n(C) + 2}$

## Heuristic Functions

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- Coverage:  $Cov = n^+(C) - n^-(C)$
- Informativity:  $Inf = \log_2(Acc)$
- Weighted relative accuracy:  $WRAcc = p(C)(p(+|C) - p(+))$



## Coverage Tests

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- How to test coverage of an example  $e$  for a clause  $C = h \leftarrow \text{Body}$ ?
- Intensional Coverage: try to derive  $e$  from  $B \cup P \cup \{C\}$
- Extensional coverage:
  - let  $\theta$  be the mgu of  $e$  and  $h$ ,
  - try to derive  $\text{Body}\theta$  from  $B \cup E^+$

## FOIL [Quinlan 90]

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- Top-down system with
  - Dimension of the beam: 1
  - Heuristic: (approximately) weighted gain of  $\text{Inf}$ :  
 $H = n(C')(Inf(C') - Inf(C))$
  - Extensional coverage
  - Refinement operator: addition of a literal or unification
  - Sufficiency criterion:  $E^+ = \emptyset$
  - Necessity criterion:  $n^-(\text{BestClause}) = 0$

## Progol [Muggleton 95]

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- Top-down system with
  - Dimension of the beam: user defined
  - Heuristic: Compression:  
 $\text{Comp} = n^+(C) - n^-(C) - |C|$
  - Intensional coverage
  - Refinement operator: see next slides
  - Sufficiency criterion:  $E^+ = \emptyset$
  - Necessity criterion:  $\text{Beam} = \emptyset$  or a maximum number of iterations of the loop is reached

## Progol Refinement Operator

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- Progol refinement operator
  - adds a literal from the most specific clause  $\perp$  after having substituted some of the constants with variables

## How to Obtain $\perp$

$B = \{parent(john, mary), male(john), parent(david, steve), male(david), parent(kathy, ellen), female(kathy)\}$   
 $e = father(john, mary)$

$$B \wedge C \models e$$

$$B \wedge \bar{e} \models \bar{C}$$

- $\bar{C}$  is a set of ground skolemized facts
- $B \wedge \bar{e} = \{parent(john, mary), male(john), parent(david, steve), male(david), parent(kathy, ellen), female(kathy), \neg father(john, mary)\}$

## How to Obtain $\perp$

$B \wedge \bar{e} = \{parent(john, mary), male(john), parent(david, steve), male(david), parent(kathy, ellen), female(kathy), \neg father(john, mary)\}$

- Let  $\bar{\perp}$  be the potentially infinite conjunction of ground literals true in the model of  $B \wedge \bar{e}$

$$\bar{\perp} = B \wedge \bar{e}$$

$\perp = father(john, mary) \leftarrow parent(john, mary), male(john), parent(david, steve), male(david), parent(kathy, ellen), female(kathy)$

## Progol Refinement Operator

- $\bar{\perp}$  is the potentially infinite conjunction of ground literals true in the model of  $B \wedge \bar{e} \Rightarrow$

$$\bar{\perp} \models \bar{C} \Rightarrow$$

$$C \models \perp$$

- We can find some of the solutions of the learning problem by looking for clauses that subsume  $\perp$
- Example:  
 $C = father(X, Y) \leftarrow parent(X, Y), male(X)$
- $C$   $\theta$ -subsumes  
 $\perp = father(john, mary) \leftarrow parent(john, mary), male(john), parent(david, steve), male(david), parent(kathy, ellen), female(kathy)$

## Progol Refinement Operator

- $\perp$  can have infinite cardinality
  - Progol uses a limit to the depth of the derivations with which  $\perp$  is built
- Moreover, restrictions to constrain the space of clauses that subsume  $\perp$ 
  - limit to the depth of variables in  $C$
  - mode declarations

## Mode Declarations

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- $modeh(n, atom)$  or  $modeb(n, atom)$
- $n$  is the maximum number of times that atom can be added to  $C$
- $atom$  is of the form  $p(Term_1, \dots, Term_n)$
- $Term_i$  is of the form
  - $+type$ : input variable of type  $type$
  - $-type$ : output variable of type  $type$
  - $\#type$ : constant of type  $type$

## Examples of Mode Declarations

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$modeh(1, plus(+int, +int, -int))$   
 $modeb(*, append(-list, +list, +list))$   
 $modeb(1, append(+list, [+any], -list))$   
 $modeb(4, (+int > \#int))$

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## Learning from Interpretations

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- Aim: learning a classifier for logical interpretations
- Classifier: a set of disjunctive clauses
- Disjunctive clause
  - $C = h_1 \vee h_2 \vee \dots \vee h_n \leftarrow b_1, b_2, \dots, b_m$
- $head(C) = \{h_1, h_2, \dots, h_n\}$
- $body(C) = \{b_1, b_2, \dots, b_m\}$
- $body^+(C) =$  set of positive literals of  $body(C)$
- $body^-(C) =$  set of atoms of negative literals of  $body(C)$
- Interpretation = set of ground atoms.

## Learning from Interpretations

- Set of clauses as a classifier
  - an interpretation is positive if all the clauses are true in the interpretation
  - an interpretation is negative if there exists at least one clause that is false in it
- A clause  $C$  is true in an interpretation  $I$  if for all grounding substitutions  $\theta$  of  $C$ :  
 $I \models body(C)\theta \rightarrow head(C)\theta \cap I \neq \emptyset$   
or  
 $body^+(C)\theta \subseteq I \wedge body^-(C)\theta \cap I = \emptyset \rightarrow head(C)\theta \cap I \neq \emptyset$

## Test of the Truth of a Clause

- Range restricted clause: all the variables in the head appear in the body
- Range restricted clause  $C$ , finite interpretation  $I$ : run the query  $? - body(C), not head(C)$  against a logic program containing  $I$
- If  $C = h_1 \vee h_2 \vee \dots \vee h_n \leftarrow b_1, b_2, \dots, b_m$  then the query is  $? - b_1, b_2, \dots, b_m, not h_1, not h_2, \dots, not h_n$
- If the query succeeds,  $C$  is false in  $I$ . If the query fails,  $C$  is true in  $I$  [De Raedt, Bruynooghe 93]

## Example

- $I = \{female(liz), male(richard), gorilla(liz), gorilla(richard)\}$
- $C = male(X) \vee female(X) \leftarrow gorilla(X)$ : the clause is true in  $I$  because the query  $? - gorilla(X), not male(X), not female(X)$  fails
- $C = male(X) \leftarrow gorilla(X)$ : the clause is false in  $I$  because the query  $? - gorilla(X), not male(X)$  succeeds with  $\theta = \{X/liz\}$ .

## Learning from Interpretations

- **Given**
  - a space of possible clausal theories  $\mathcal{H}$
  - a set  $P$  of interpretations
  - a set  $N$  of interpretations
- **Find**: a clausal theory  $H \in \mathcal{H}$  such that
  - for all  $p \in P, p \models H$
  - for all  $n \in N, n \not\models H$
- Less expressive than learning from entailment: no recursive definitions

## Test with Background

- Background: a normal program  $B$
- Truth of a clause  $C$  in the interpretation  $M(B \cup I)$  where  $M$  is the model according to the chosen semantics and  $I$  is an interpretation (i.e.  $B \cup I \models C$ )
- Range restricted clause  $C$ , normal program  $B$  containing only range restricted clauses, interpretation  $I$ : run the query  $? - body(C), not head(C)$  against the logic program  $B \cup I$ .
- If the query succeeds,  $C$  is false in  $M(B \cup I)$  ( $B \cup I \not\models C$ ). If the query fails,  $C$  is true in  $M(B \cup I)$  ( $B \cup I \models C$ )

## Learning from Int. with Background

### Given

- a space of possible clausal theories  $\mathcal{H}$
- a set  $P$  of interpretations
- a set  $N$  of interpretations
- a background theory  $B$

**Find:** a clausal theory  $H \in \mathcal{H}$  such that

- for all  $p \in P$ ,  $B \cup p \models H$
- for all  $n \in N$ ,  $B \cup n \not\models H$

## Generality Relation

- $cover(\{C\}, e) = true$  if  $e \models C$
- $C \geq D \Rightarrow C \models D \Rightarrow D$  is more general than  $C$
- the relation is reversed
- Example:

```
false ← true
false ← gorilla(X)
female(X) ← gorilla(X)
female(X) ∨ male(X) ← gorilla(X)
```

## ICL [De Raedt, Van Laer, 95]

- Dual version of a top down entailment algorithm:
  - coverage loop is performed on negative examples
- Updates CN2 to first order

**ICL**( $P, N, B$ )

$H := \emptyset$

repeat

$C := \text{FindBestClause}(P, N, B)$

    if  $C \neq null$  then

        add  $C$  to  $H$

        remove from  $N$  all interpretations that are false for  $C$

until  $C = null$  or  $N$  is empty

return  $H$

## ICL FindBestClause

```
FindBestClause( $P, N, B$ )
  Beam := {false ← true}, BestClause := null
  while Beam is not empty do
    NewBeam := ∅
    for each clause  $C$  in Beam do
      for each refinement  $Ref$  of  $C$  do
        if  $Ref$  is better than  $BestClause$  and  $Ref$  is
           statistically significant then
          BestClause :=  $Ref$ 
        if  $Ref$  is not to be pruned then
          add  $Ref$  to  $NewBeam$ 
          if size of  $NewBeam$  >  $MaxBeamSize$  then
            remove worst clause from  $NewBeam$ 
    Beam :=  $NewBeam$ 
  return BestClause
```

Inductive Logic Programming – p. 53/93

## ICL Heuristics

- $n(\bar{C})$  = number of interpretations (positive and negative) where  $C$  is false
- $n^-(\bar{C})$  = number of negative interpretation where  $C$  is false
- $H(C) = p(-|\bar{C}) = \frac{n^-(\bar{C})+1}{n(\bar{C})+2}$  = precision over negative class

Inductive Logic Programming – p. 54/93

## Tilde [Blockeel, De Raedt 98]

- Updates C4.5 to first-order
- Solve a slightly different learning problem
- **Given**
  - a space of possible theories  $\mathcal{H}$
  - a set of classes  $C$
  - a set  $E$  of classified interpretations (couples  $(I, c)$ )
  - a background theory  $B$
- **Find:** a theory  $H \in \mathcal{H}$  such that  $\forall (I, c) \in E$ 
  - $B \cup I \cup H \models c$
  - $\forall c' \in C \setminus \{c\}, : B \cup I \cup H \not\models c'$
- Less expressive than learning from entailment: excludes recursive definitions

Inductive Logic Programming – p. 55/93

## Example

```
 $E =$ 
Machine 1: {worn(gear), worn(engine), sendback}
Machine 2: {ok}
Machine 3: {worn(gear), fix}
Machine 4: {worn(engine), sendback}
Machine 5: {worn(gear), worn(chain), fix}
 $B =$ 
replaceable(gear).
replaceable(wheel).
replaceable(chain).
not_replaceable(engine).
not_replaceable(control_unit).
```

Inductive Logic Programming – p. 56/93

## Example

---

```
worn(A) ?
+--yes: not_replaceable(A) ?
|      +--yes: [sendback] [6.0/6.0]
|      +--no:  [fix] [6.0/6.0]
+--no:  [ok] [3.0/3.0]
```

Equivalent Prolog program (FODL):

```
class(sendback) :- worn(A),not_replaceable(A), !.
class(fix) :- worn(A), !.
class(ok).
```

## Tilde Algorithm

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**Tilde**( $E$ )

$T := \text{GrowTree}(E, \text{true})$

return Prune( $T$ )

**GrowTree**( $E, Q$ : query)

$Q_b := \text{OptimalSplit}(\rho(Q), E)$

if StopCrit( $Q_b, E$ ) then

    return *leaf*(majority\_class( $E$ ))

else

$\text{conj} := Q_b - Q$

$E_1 := \{I \in E \mid \leftarrow Q_b \text{ succeeds in } B \cup I\}$

$E_2 := \{I \in E \mid \leftarrow Q_b \text{ fails in } B \cup I\}$

$T := \text{inode}(\text{conj},$

        GrowTree( $E_1, Q_b$ ), GrowTree( $E_2, Q$ ))

    return  $T$

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## Descriptive ILP

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- Discovering regularities, patterns
- Example tasks:
  - finding association rules
  - clustering
  - subgroup discovery

## Claudien [De Raedt, Dehaspe 97]

- Learning problem: Given
  - a space of possible clausal theories  $\mathcal{H}$
  - a set  $P$  of interpretations
  - a background theory  $B$
- **Find:** a clausal theory  $H \in \mathcal{H}$  such that
  - $\forall p \in PB \cup p \models H$
  - $H$  is maximally specific

## Example

$p1 = \{female(liz), male(richard),$   
 $gorilla(liz), gorilla(richard)\}$   
 $p2 = \{female(ginger), male(fred),$   
 $gorilla(ginger), gorilla(fred)\}$

If  $\mathcal{H}$  is restricted to range-restricted, constant-free clauses a solution is:

$gorilla(X) \leftarrow female(X)$   
 $gorilla(X) \leftarrow male(X)$   
 $male(X) \vee female(X)$   
 $\leftarrow male(X), female(X)$

## Claudien Algorithm

**ClausalDiscovery**( $E, B$ )

$H := \emptyset$

$Beam := \{false \leftarrow true\}$

while  $Beam$  is not empty do

  delete from  $Beam$  the first clause  $C$

  if  $C$  is true on  $E$  then

$H := H \cup \{C\}$

  else

    for all  $C' \in \rho(C)$  for which not  $prune(C')$  do

$Beam := Beam \cup \{C'\}$

return  $H$

## Outline of the Talk

- Predictive ILP
  - Learning from entailment
    - Bottom-up systems: Golem
    - Top-down systems: FOIL, Progol
  - Learning from interpretations
    - ICL, Tilde
- Descriptive ILP
  - Claudien
- Probabilistic ILP
  - ALLPAD
- Applications



## Probabilistic ILP

- Learning a probabilistic logic program
- Formalisms:
  - BLP: Bayesian Logic Programs
  - CLP(BN): CLP over Bayesian Networks
  - MIA: Meta-Interpreter Approach
  - SLP: Stochastic Logic Programs
  - LBN: Logical Bayesian Networks
  - LPAD: Logic Programs with Annotated Disjunctions

## LPADs [Vennekens et al. 04]

- A **Logic Program with Annotated Disjunctions** consists of a set of formulas of the form

$$(h_1 : p_1) \vee (h_2 : p_2) \vee \dots \vee (h_n : p_n) \leftarrow b_1, b_2, \dots, b_m$$

- The  $p_i$  are real numbers in the interval  $[0, 1]$  such that  $\sum_{i=1}^n p_i = 1$ .
- Instance: a ground normal program obtained by selecting one atom from each grounded clause
- They define a probability distribution  $\pi_P$  over
  - instances: product of probabilities of heads selected
  - interpretations and formulas: probability of  $\phi$  is the sum of the probabilities of instances where  $\phi$  is true

## Example

*mother(m, c) father(f, c)*

*cg(m, 1, w) cg(m, 2, w)*

*cg(f, 1, p) cg(f, 2, w)*

$(cg(X, 1, A) : 0.5) \vee (cg(X, 1, B) : 0.5) \leftarrow$   
*mother(Y, X), cg(Y, 1, A), cg(Y, 2, B)*

$(cg(X, 2, A) : 0.5) \vee (cg(X, 2, B) : 0.5) \leftarrow$   
*father(Y, X), cg(Y, 1, A), cg(Y, 2, B)*

*color(X, purple) ← cg(X, \_N, p)*

*color(X, white) ← cg(X, 1, w), cg(X, 2, w)*

## Learning LPADs

- Learning problem:
  - Given:**
    - a set  $E$  of couples  $(I, \pi(I))$  where  $I$  is an interpretation and  $\pi(I)$  is its associated probability
    - a space of possible LPAD  $\mathcal{H}$
  - Find:**
    - an LPAD  $P \in \mathcal{H}$  such that  $\forall (I, \pi(I)) \in E \pi_P^*(I) = \pi(I)$
- Instead of a set of couples  $(I, \pi(I))$ , the input of the learning problem can be a multiset  $E'$  of interpretations.

## ALLPAD [Riguzzi 06]

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- Three phases:
  - find all the ground clauses that satisfy a number of constraints
  - compute the probabilities in the head
  - solve an optimization problem

## Outline of the Talk

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- Predictive ILP
  - Learning from entailment
    - Bottom-up systems: Golem
    - Top-down systems: FOIL, Progol
  - Learning from interpretations
    - ICL, Tilde
- Descriptive ILP
  - Claudien
- Probabilistic ILP
  - ALLPAD
- **Applications**

## Applications

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- Biology
- Chemistry
- Engineering
- Various

## Algorithm Evaluation

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- Notation:
  - $n^+(P)$  number of positive examples covered by  $P$
  - $n^-(\bar{P})$  number of negative examples not covered by  $P$
  - $n = |E|$
- Accuracy:

$$Acc(P) = \frac{n^+(P) + n^-(\bar{P})}{n}$$

## Structure Activity Relationships (SARs)

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- Predicting the activity of a compound on humans based on its chemical structure and properties
  - Drugs: whether they are effective
  - Compounds, drugs: whether they are toxic

## Description of Chemical Compounds

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### Basic structure:

*atom(compound, atom, element, atomType, charge)*

e.g. *atom(d2, d2\_1, c, 22, 0.067)*

*bond(compound, atom1, atom2, bondType)*

e.g. *bond(d2, d2\_1, d2\_2, 7)*

### Structures:

*benzene(compound, listOfAtoms)*

e.g. *benzene(d4, [d4\_6, d4\_1, d4\_2, d4\_3, d4\_4, d4\_5])*

*phenanthrene(compound, listOfListsOfAtoms)*

*nitro(compound, listOfAtoms)*

...

### Properties:

*polar(atom, polarity)*

*polar(d2\_1, polar3)*

...

## SAR

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- Drugs against Alzheimer's disease
  - Golem: not significantly different from propositional, comprehensibility [King et al. 95]
- Drugs for inhibition of E. Coli Dihydrofolate Reductase
  - Golem: not significantly different from propositional, comprehensibility [King et al. 95]
- Predicting carcinogenicity
  - Progol: 72% highest machine accuracy [Srinivasan et al. 97]

## SAR

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- Predicting mutagenicity
  - regression friendly compounds
    - FOIL: 82% [Srinivasan et al 95]
    - ICL: 86.2% [Van Laer et al. 97]
    - Progol: 88% [Srinivasan et al 95]
    - Claudien: found alternative explanations [De Raedt, Dehaspe 97]
  - regression unfriendly compounds
    - Progol: 85.7% [King et al. 96]

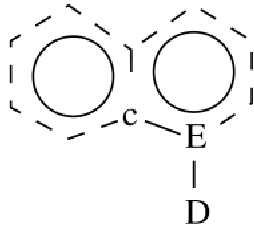
# Progol on Mutagenesis

$active(A) \leftarrow$   
 $atom(A, B, c, 27, C),$   
 $bond(A, D, E, 1), bond(A, E, B, 7)$

A carbon atom of type 27 merges two six-membered aromatic rings.

A bond of type 7 is an aromatic bond.

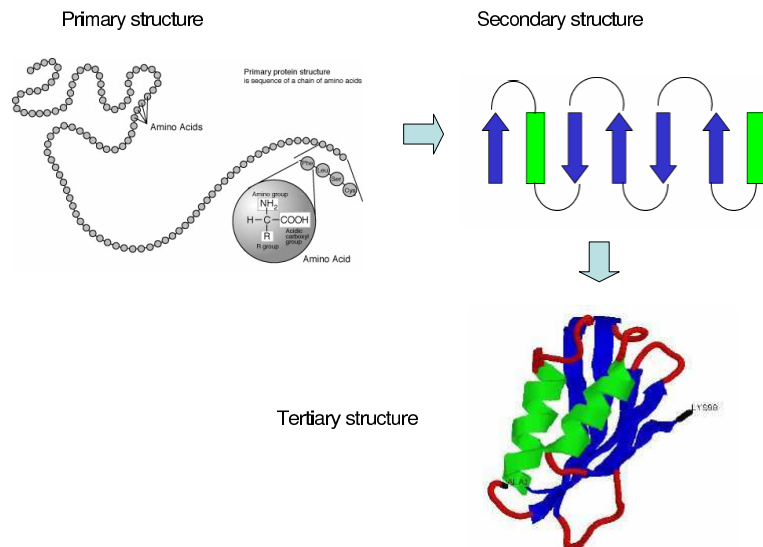
This rule identifies compounds of two fused six-membered aromatic rings, one of which has a further single bond with an atom of any type.



# Biology

- Description of the binding sites (pharmacophores) of ACE inhibitors (hypertension drug) and an HIV-protease inhibitor (an anti-AIDS drug)
  - Progol: rediscovered a pharmacophore found by experts [Finn et al. 98]
- Biological classification of river water quality
  - Golem: comprehensibility [Dzeroski et al. 94]
  - Claudien: intuitive rules [De Raedt, Dehaspe 97]

# Proteins



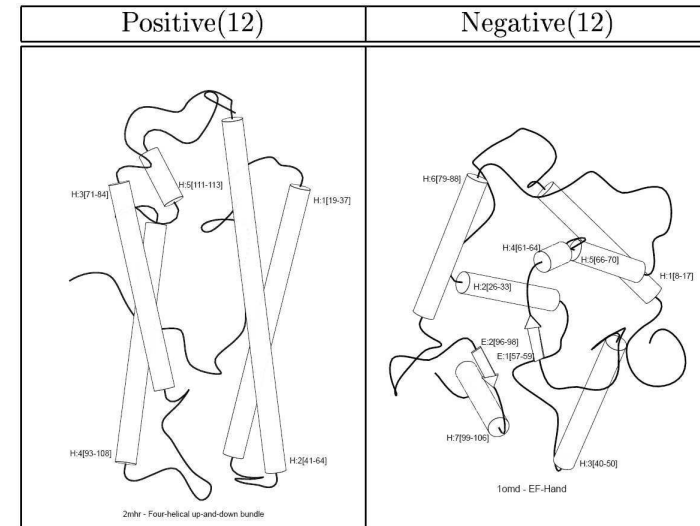
# Protein Secondary Structure

- Predicting protein secondary structure from the amino-acid sequence
- Structures
  - helices, of various types and length
  - strands, of various orientations and length
- Results:
  - Golem: 80% [Muggleton et al. 92]
  - FOIL: 65% [Quinlan, Cameron-Jones 95]

# Protein Tertiary Structure

- Predicting the tertiary structure of proteins by classifying them into one of the SCOP classes
- Proteins represented as a sequence of secondary structure elements
- Results:
  - Progol: 78.28% [Turcotte et al. 01]
  - ALLPAD: 85.67% [Riguzzi 06]

# Protein Tertiary Structure



# Chemistry

- Identification of the structure of diterpene from spectral information
  - FOIL: 78.3% [Dzeroski et al. 96]
  - Tilde: 90.4% [Blockeel, De Raedt 98]
- Predicting the half-time of aqueous biodegradation of a compound from its chemical structure
  - ICL: 58.1% [Van Laer et al. 97]
- Judging whether a molecule is a musk
  - Tilde: 79.4% [Blockeel, De Raedt 98]

# Engineering

- Learning rules for finite element mesh design
  - Claudien: 34% from pos. ex. only [De Raedt, Dehaspe 97]
  - ICL: 66.5% [Van Laer et al. 97]
  - Golem: 78% [Dolsak et al. 94]

## Various

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- Identifying document components
  - FOIL: 96.3%-100% [Quinlan, Cameron-Jones 95]
- Recovering program loop invariants from program traces
  - Claudien: found true invariants [De Raedt, Dehaspe 97]

## Conclusions

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- Active Research Field
- New directions:
  - upgrading propositional systems (probabilistic models, support vector machines, reinforcement learning, neural networks,...)
  - downgrading first order systems (to make it more efficient in special cases)
  - multi relational data mining

## Pointers

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- ILPnet2
  - <http://www.cs.bris.ac.uk/~ILPnet2/>
  - <http://www-ai.ijs.si/~ilpnet2/>
- KDnet <http://www.kdnet.org/>
- Books:
  - [Lavrac, Dzeroski 94]: freely available in pdf on the web
  - [Bergadano et al. 96]
  - [Dzeroski, Lavrac 01]

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